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Worksheet 6

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Fundamental Algorithms

Exercise 1

– obsolete –

Exercise 2

Consider a partitioning algorithm that, in the worst case, will partition an array of m elements into two partitions of size $\lfloor \epsilon m \rfloor$ and $\lceil (1-\epsilon)m \rceil$, where ϵ is fixed, and $0 < \epsilon < 1$. Show that a quicksort algorithm based on this partitioning has a worst-case complexity of $O(n \log n)$.