

# Algorithms of Scientific Computing

## Exercise 1: Fast Discrete Cosine Transform

The butterfly scheme is retrieved as usual:

$$\begin{aligned}
 F_k &= \frac{1}{2N} \sum_{n=-N+1}^N f_n \omega_{2N}^{-kn} = \frac{1}{2} \left( \frac{1}{N} \sum_{n=-N/2+1}^{N/2} f_{2n} \omega_{2N}^{-2kn} + \frac{1}{N} \sum_{n=-N/2+1}^{N/2} f_{2n-1} \omega_{2N}^{-k(2n-1)} \right) \\
 &= \frac{1}{2} \left( \underbrace{\frac{1}{N} \sum_{n=-N/2+1}^{N/2} f_{2n} \omega_N^{-kn}}_{=:G_k} + \underbrace{\frac{1}{N} \sum_{n=-N/2+1}^{N/2} f_{2n-1} \omega_N^{-kn} \omega_{2N}^k}_{=:H_k} \right) \\
 &= \frac{1}{2} \left( G_k + \omega_{2N}^k H_k \right) \\
 F_{k+N} &= \frac{1}{2} \left( G_{k+N} + \omega_{2N}^{k+N} H_{k+N} \right) = \frac{1}{2} \left( G_k - \omega_{2N}^k H_k \right)
 \end{aligned}$$

For the datasets  $g_n := f_{2n}$  and  $h_n := f_{2n-1}$ , respectively, we can try to find other symmetries:

$$g_{-n} = f_{2(-n)} = f_{-2n} = f_{2n} = g_n$$

The "even" data also shows an even symmetry and therefore leads to another Cosine Transform but with half length.

Analogously, for the data with odd indices:

$$h_{-n} = f_{2(-n)-1} = f_{-2n-1} = f_{2n+1} = f_{2(n+1)-1} = h_{n+1}$$

Again we get an "even" symmetry. This is the transform shown in the lecture, known as Quarter-Wave-DCT, again with half length.

## Exercise 2: DFT and Least Squares Approximation

Since the Euclidean norm is defined by a scalar product, the function giving the error  $E(\alpha_0, \dots, \alpha_{N-1}) = \|\mathbf{f} - \Phi_N(x)\|_2^2$  is a quadratic function that attains its extremum where  $\nabla E(\alpha_0, \dots, \alpha_{N-1})$  vanishes.

Given that the  $k$ -th partial derivative is given by

$$\frac{\partial E}{\partial a_k} = \sum_{n=0}^{N-1} \left[ e^{-i2\pi nk/N} \left( f_n - \sum_{p=0}^{N-1} \alpha_p e^{i2\pi np/N} \right) \right], \quad (1)$$

we set each partial derivative to zero and obtain

$$\sum_{n=0}^{N-1} \left[ e^{-i2\pi nk/N} \left( f_n - \sum_{p=0}^{N-1} \alpha_p e^{i2\pi np/N} \right) \right] = 0. \quad (2)$$

Rearranging the terms gives us the set of  $N$  equations

$$\sum_{n=0}^{N-1} f_n \omega_N^{-nk} = \sum_{p=0}^{N-1} \alpha_p \underbrace{\sum_{n=0}^{N-1} \omega_N^{n(p-k)}}_{N\delta_N(p-k)}, \quad (3)$$

where  $\omega_N = e^{i2\pi/N}$  and  $\delta_N(k)$  is known as the modular Kronecker delta defined by

$$\delta_N(k) = \begin{cases} 1 & \text{if } k = 0 \text{ or a multiple of } N, \\ 0 & \text{otherwise.} \end{cases} \quad (4)$$

Next, we will show that  $\sum_{n=0}^{N-1} \omega_N^{n(p-k)} = N\delta_N(p-k)$ .  $\omega_N^k$  are complex exponentials that for  $k = 0, \dots, N-1$  are the  $N$ th roots of unity.

$$\underbrace{(\omega_N^k)^N}_{:=z} = \left( e^{i2\pi k/N} \right)^N = e^{i2\pi k} = 1 \quad (5)$$

We can rewrite  $z^N - 1 = 0$  as

$$z^N - 1 = (z - 1) \left( z^{N-1} + z^{N-2} + \dots + z + 1 \right) = (z - 1) \sum_{n=0}^{N-1} z^n = 0. \quad (6)$$

A product is zero, if one of the factors evaluates zero. We have a case distinction.

1.  $z = \omega_N^{p-k}$  where  $p-k$  is not a multiple of  $N$ . Then  $z \neq 1$ . The product can only evaluate zero if the second term evaluates zero. Hence

$$\sum_{n=0}^{N-1} z^n = \sum_{n=0}^{N-1} \left( \omega_N^{p-k} \right)^n = 0 \quad (7)$$

2.  $z = \omega_N^{p-k}$  where  $p-k$  is a multiple of  $N$ . Then  $\omega_N^{p-k} = 1$  and

$$\sum_{n=0}^{N-1} z^n = \sum_{n=0}^{N-1} \left( \omega_N^{p-k} \right)^n = \sum_{n=0}^{N-1} 1 = N. \quad (8)$$

Using this result in equation (3) we obtain

$$\sum_{n=0}^{N-1} f_n \omega_N^{-nk} = N\alpha_k, \quad (9)$$

for  $k = 0, \dots, N-1$ . We get that  $\alpha_k$ s correspond exactly to the DFT coefficients

$$\alpha_k = \frac{1}{N} \sum_{n=0}^{N-1} f_n \omega_N^{-nk}. \quad (10)$$

The least squares error is minimised when the coefficients of the trigonometric polynomial correspond to the DFT coefficients.

The FFT algorithm can be used to efficiently compute the  $\alpha_k$ .

### Exercise 3: Fast Discrete Sine Transform

The butterfly scheme is retrieved as usual:

$$\begin{aligned} F_k &= \frac{1}{2N} \sum_{n=-N+1}^N f_n \omega_{2N}^{-kn} = \frac{1}{2} \left( \frac{1}{N} \sum_{n=-N/2+1}^{N/2} f_{2n} \omega_{2N}^{-2kn} + \frac{1}{N} \sum_{n=-N/2+1}^{N/2} f_{2n-1} \omega_{2N}^{-k(2n-1)} \right) \\ &= \frac{1}{2} \left( \underbrace{\frac{1}{N} \sum_{n=-N/2+1}^{N/2} f_{2n} \omega_N^{-kn}}_{=:G_k} + \underbrace{\frac{1}{N} \sum_{n=-N/2+1}^{N/2} f_{2n-1} \omega_N^{-kn} \omega_{2N}^k}_{=:H_k} \right) \\ &= \frac{1}{2} \left( G_k + \omega_{2N}^k H_k \right) \\ F_{k+N} &= \frac{1}{2} \left( G_{k+N} + \omega_{2N}^{k+N} H_{k+N} \right) = \frac{1}{2} \left( G_k - \omega_{2N}^k H_k \right) \end{aligned}$$

For the datasets  $g_n := f_{2n}$  and  $h_n := f_{2n-1}$ , respectively, we can try to find other symmetries:

$$g_{-n} = f_{2(-n)} = -f_{-2n} = -f_{2n} = -g_n$$

The "odd" data also shows an odd symmetry and therefore lead to another Sine Transform but with half length.

Analog for the data with odd indices:

$$h_{-n} = f_{2(-n)-1} = f_{-2n-1} = -f_{2n+1} = -f_{2(n+1)-1} = -h_{n+1}$$

Again we get an "odd" symmetry. However, this is the transform shown in the lecture, known as Quarter-Wave-DST, again with half length.

For a dataset with the symmetry constraint  $f_{-n} = -f_{n+1}$  we get accordingly

$$g_{-n} = f_{2(-n)} = f_{-2n} = -f_{2n+1} = -h_{n+1}$$

and

$$h_{-n} = f_{-2n-1} = f_{-2n+1} = -f_{2n+2} = -f_{2n-1} = -g_{n+1}$$