

## Parallel Numerics

### Exercise 11: Previous Exam Questions

## 1 Preconditioning & Iterative Solvers

(From 2016)

Given is the system of linear equations  $Ax = b$  with

$$A = \begin{pmatrix} 1 & 0 & 3 \\ 2 & 1 & 0 \\ 0 & -1 & -1 \end{pmatrix}, b = \begin{pmatrix} 1 \\ 2 \\ 1 \end{pmatrix}.$$

The exact solution of this system is  $x = \left(\frac{10}{7}, -\frac{6}{7}, -\frac{1}{7}\right)^T$  and does not have to be calculated separately.

- a) Apply the static SPAI algorithm to determine a part of preconditioner  $\tilde{M}$  for matrix  $A$  with sparsity pattern  $\mathcal{J}_1 = \{3\}$ . ( $\approx 3$  credits)

**Hint:** You can use the *normal equations* to solve a least squares problem.

- b) Combine your result of subtask a) with  $\begin{pmatrix} 0 & 0.4 & 0 \\ 0 & 0 & -0.5 \\ * & 0 & 0 \end{pmatrix}$  to get an actual right-preconditioner

$M$ . Precondition  $Ax = b$  with  $M$  and calculate two steps of the Jacobi iteration of the preconditioned system. Use  $y^{(0)} = (0, 0, 0)^T$  as start vector. ( $\approx 4$  credits)

(If you did not get any result in subtask a), then use  $\tilde{M}_1 = (0, 0, 0.6)^T$ .)

- c) Without preconditioning, Jacobi gives the following sequence of solutions  $x^{(i)}$  for  $Ax = b$  (again with a start vector  $x^{(0)} = (0, 0, 0)^T$ ):

$$x^{(1)} = \begin{pmatrix} 1 \\ 2 \\ -1 \end{pmatrix}, x^{(2)} = \begin{pmatrix} 4 \\ 0 \\ -3 \end{pmatrix}, x^{(3)} = \begin{pmatrix} 10 \\ -6 \\ -1 \end{pmatrix}, \dots, x^{(10)} = \begin{pmatrix} 94 \\ -618 \\ 185 \end{pmatrix}, \dots$$

Compare this result with the actual solution  $x$  and your solution for the preconditioned system from subtask b). What's the reason for your observation? ( $\approx 1$  credit)

(If you did not get any result in subtask b), then just compare the result from above with the actual solution  $x$ .)

## 2 Sparse Gauss Elimination

(From 2015)

Given is the sparse symmetric  $5 \times 5$  matrix

$$A = \begin{pmatrix} 4 & 1 & 2 & \frac{1}{2} & 2 \\ 1 & \frac{1}{2} & 0 & 0 & 0 \\ 2 & 0 & 3 & 0 & 0 \\ \frac{1}{2} & 0 & 0 & \frac{5}{8} & 0 \\ 2 & 0 & 0 & 0 & 16 \end{pmatrix}.$$

- Let  $A$  be an adjacency matrix. Draw the corresponding graph  $G(A)$ ! ( $\approx 1$  credit)
- Use Gaussian elimination without any pivoting to eliminate the first column of  $A$ ! What can be observed after this elimination regarding the remaining matrix elements in comparison to the original matrix  $A$ ? ( $\approx 2$  credits)
- The *multiple minimum degree reordering* is one way to do pivoting for sparse matrices. It works as follows:
  - Define  $r_m :=$  number of nonzero entries in row  $m$ .
  - Choose pivot index  $i$  by  $r_i = \min_m r_m$ . If this holds for multiple  $i$ , then chose one of these  $i$  randomly.
  - Do the pivot permutation and the elimination.

Use Gaussian elimination with *multiple minimum degree reordering* to eliminate the first column of  $A$ ! What can be observed after this elimination regarding the remaining matrix elements in comparison to the result of subtask b)? ( $\approx 3$  credits)

**Hint:** Since  $A$  is a symmetric matrix, rows **and** columns have to be interchanged when pivoting to keep the symmetric structure.

- Briefly describe how *multiple minimum degree reordering* can be used to parallelize Gaussian elimination for sparse matrices! ( $\approx 1$  credit)

## 3 Parallel Jacobi

(From 2013)

From the lecture and the tutorials you are familiar with the Jacobi stationary iterative method for solving a linear system  $Ax = b$  where  $A$  is an  $n \times n$  matrix,  $b$  the given right-hand side vector and  $x$  the solution vector to be computed.

Give a parallel implementation of Jacobi's method by answering the following questions. Assume that a block distribution of all vectors and a block-row distribution of  $A$  is available on all processing elements (PEs). The following variables used in the source code are placeholders for:

**p**                    number of processing elements (PEs)  
**n**                    dimension of the underlying problem  
**my\_rank**            rank ID of current PE  
**max\_iter**            maximum number of iterations to perform  
**A\_local**            local block-row distribution of  $A$   
**b\_local, x\_local**    local block distribution of  $b$  and  $x$ , respectively

- a) Several data has to be available on all processing elements (PEs). Give the MPI/C-code to make the dimension of the underlying problem (variable **n**) and the maximum number of iterations (variable **max\_iter**) available on all PEs. Provide only the requested source code lines, i.e. no complete program is necessary.
- b) Formulate the component-wise representation of the Jacobi method for  $x_j^{(k+1)}$ . Give additionally a **pseudocode** formulation to compute the complete vector  $x^{(k+1)}$ .
- c) Complete the following routine `void Parallel_Jacobi(...)` with **own source code at line 22**. The routine should compute a parallel version of the stationary Jacobi method. Assume that **n** is a multiple of **p** and that the macro `Swap(x,y)` in line **1** is available. Note that the iteration is stopped if **max\_iter** is reached or  $\|x_{\text{new}} - x_{\text{old}}\|_2 < 10^{-6}$ .

```

1:#define Swap(x,y) {float* temp; temp = x; x = y; y = temp;}
2:
3:void Parallel_Jacobi(MATRIX_T A_local, float x_local[], float b_local[],
4:                    int n, int max_iter, int p, int my_rank) {
5:    int i_local, i_diag, j, iter_num;
6:    float x_temp1[MAX_DIM], x_temp2[MAX_DIM];
7:    float* x_old;
8:    float* x_new;
9:
10:   /* Initialize local temporary x */
11:   MPI_Allgather(b_local, n/p, MPI_FLOAT, x_temp1, n/p, MPI_FLOAT,
12:               MPI_COMM_WORLD);
13:
14:   x_new = x_temp1;
15:   x_old = x_temp2;
16:   iter_num = 0;
17:   do {
18:       iter_num++;
19:       /* Interchange x_old and x_new */
20:       Swap(x_old, x_new);
21:
22:       /* *** insert code for parallel Jacobi here *** */
23:
24:   } while ((iter_num < max_iter) && (norm(x_new,x_old,n) >= 1e-6));
25:}
  
```