

# A Look-ahead Algorithm for Solving Nonsymmetric Linear Toeplitz Equations

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## Abstract

Fast solvers for linear Toeplitz equations like the Levinson algorithm are known to be stable only for symmetric positive definite matrices. In general they may break down caused by a division with zero in view of a singular submatrix. One way to overcome this breakdown uses the connection between Toeplitz matrices and orthogonal polynomials. Here, we will present another approach for dealing this case. In the case of a breakdown we introduce perturbations of the original matrix in order to make the submatrix regular. Then we can step forward without breakdown, and after a few steps we have to return to the original matrix.

## 1 Look-ahead Algorithm for Computing Gohberg-Semencul Formulas

The solution of linear equations  $T_n x = b$  with  $T_n = (t_{i-j})_{i,j=0}^n$  a Toeplitz matrix appears in many applications (see e.g. [2]). If we want to solve more than one Toeplitz equation then usually we solve one or two special linear systems in order to determine a Gohberg-Semencul formula that expresses  $T_n^{-1}$  as the sum of products of upper and lower Toeplitz matrices. Given a Gohberg-Semencul formula we can solve every equation  $T_n x = b$  in  $O(n \log(n))$  steps. For the positive definite case the original Gohberg-Semencul formula [1,7] that we get by solving  $T_n x = e_0$  and  $T_n y = e_n$  exists,

$$T_n^{-1} = (1/x_0) \left( L_n(x)U_n(y) - L_n(ZJy)U_n(ZJx) \right)$$

but in the general case it breaks down for  $x_0 = 0$ ,  $x_0$  the first component of the vector  $x$ . The second Gohberg-Semencul formula is given by [6,9]

$$T_n^{-1} = \left( L_n(z)U_n(ZJx) - L_n(x)U_n(-e_0 + ZJz) \right)$$

with  $T_n z = (0, t_{-n}, \dots, t_{-1})^T$ ,  $e_0 = (1, 0, \dots, 0)^T$ ,  $Z$  is the lower shift matrix,  $J$  is the counteridentity  $J_{j,k} = \delta_{n+1-j,k}$ ,  $j, k = 0, \dots, n$ ,  $L_n$  denotes a lower Toeplitz matrix and  $U_n$  an upper Toeplitz matrix. This formula is always well defined. All possible Gohberg-Semencul formulas are described in [10] using the displacement representation of a Toeplitz matrix [11]. If we have computed one Gohberg-Semencul formula then we can easily get the generating vectors of all other formulas [10].

The aim of the presented algorithm is to compute generating vectors for Gohberg-Semencul formulas for a given regular Toeplitz matrix  $T_n$  recursively as fast and stable as possible. The method is based on the recurrence between the generating vectors  $x$ ,  $y$ ,

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and  $z$  that is described in [9,p.56]. Let us first recall these equations for the case that all submatrices of  $T_n$  are regular. Set

$$T_k x^k = e_0, \quad T_k z^k = \tau_-^k = (t_{-k-1} \dots t_{-1})^T, \quad T_k^T y^k = e_0, \quad T_k^T w^k = \tau_+^k = (t_{k+1} \dots t_1)^T.$$

Then we get the two different recursions

$$(1) \quad x^{k+1} = \frac{1}{1 - w_0^k z_0^k} \left( \begin{pmatrix} x^k \\ 0 \end{pmatrix} - w_0^k \begin{pmatrix} 0 \\ y^k \end{pmatrix} \right)$$

$$(2) \quad y^{k+1} = \frac{1}{1 - w_0^k z_0^k} \left( \begin{pmatrix} 0 \\ y^k \end{pmatrix} - z_0^k \begin{pmatrix} x^k \\ 0 \end{pmatrix} \right)$$

or

$$(3) \quad x^{k+1} = \begin{pmatrix} x^k \\ 0 \end{pmatrix} - \frac{w_0^k}{\tau_+^k z^k - t_0} \begin{pmatrix} z^k \\ -1 \end{pmatrix}$$

$$(4) \quad z^{k+1} = \begin{pmatrix} 0 \\ z^k \end{pmatrix} - (\tau_-^k z^k - t_{-k-2}) x^{k+1}$$

with

$$(5) \quad w_0^k = \tau_+^k x^k, \quad z_0^k = \tau_-^k x^k, \quad w_z^k = \tau_+^k z^k = \tau_-^k w^k = z_w^k, \quad \gamma^k = w_z^k - t_0, \quad z_z^k = \tau_-^k z^k.$$

The  $(x, y)$ -based formulas (1-2) have the advantage that they need only two inner products for  $w_0^k$  and  $z_0^k$  but the disadvantage that a breakdown occurs if  $T_{k-1}$  or  $T_{k+1}$  is singular. The recurrence (3-4) takes three inner products for  $w_0^k, w_z^k,$  and  $z_z^k$  but breaks down only if  $T_{k+1}$  is singular. Thus we are interested in mainly using (1-2), but we have to switch between both recurrences sometimes in order to avoid unnecessary breakdowns. To describe these switches we can use the formulas (see [10] and [9])

$$(6) \quad z^k = \frac{1}{x_0^k} (z_0^k x^k - ZJy^k) \quad \text{and} \quad y^{k+1} = \frac{1}{w_z^k - t_0} J \begin{pmatrix} z^k \\ -1 \end{pmatrix}.$$

Furthermore, we see from [10]

$$-e_0 + ZJz^k = (w_z^k - t_0)y^k - z_0^k w^k,$$

that

$$(7) \quad -(w_z^k - t_0)x_0^k = 1 - z_0^k w_0^k$$

Therefore, (1-2) breaks down if  $x_0^k = 0$  or if (3-4) breaks down. Especially for  $w_z^k - t_0 = 0$  there holds  $1 = z_0^k w_0^k$  and hence  $z_0^k \neq 0$ . Thus, if in (3-4) a breakdown occurs, that means  $w_z^k - t_0 = 0$ , then  $z_0^k$  is not equal 0 and we can force  $\tilde{w}_z^k - t_0 \neq 0$  by changing  $t_{k+1}$  to  $\tilde{t}_{k+1}$ . In the same way, if (1-2) breaks down with  $x_0^k \neq 0$  then also  $z_0 \neq 0$  and in view of (7) we can force  $1 - w_0^k z_0^k$  to be not equal to zero. Hence, we get regular extensions  $\tilde{T}_{k+1}, \dots, \tilde{T}_{k+r}$  of  $T_k$ .

In every step we now have to check whether the original matrix  $T_{k+r}$  is regular, and if this is true we have to switch from the generated vectors  $\tilde{x}^{k+r}$  and  $\tilde{y}^{k+r}$  to the original vectors  $x^{k+r}$  and  $z^{k+r}$ . Note, that we can not compute  $y^{k+r}$  because  $T_{k+r-1}$  will be ill-



conditioned. From  $(x^{k+r}, z^{k+r})$  we can compute  $(x^{k+r+1}, y^{k+r+1})$ , or in the breakdown case  $(\tilde{x}^{k+r+1}, \tilde{y}^{k+r+1})$ . Thus, the main recurrence is of the form (1-2), and only after a breakdown we use (3) and the second equation in (6).

To get  $(x^{k+r}, z^{k+r})$  we apply the Sherman-Morrison-Woodbury formula

$$(8) \quad T_{k+r}^{-1} = (\tilde{T}_{k+r} + \tilde{\Delta})^{-1} = \tilde{T}_{k+r}^{-1} - \tilde{T}_{k+r}^{-1} \begin{pmatrix} 0 \\ I_r \end{pmatrix} G_{k+r}^{-1} (\Delta_r \quad 0) \tilde{T}_{k+r}^{-1}$$

with

$$(9) \quad G_{k+r} = I_r + \Delta_r (I_r \quad 0) \tilde{T}_{k+r}^{-1} \begin{pmatrix} 0 \\ I_r \end{pmatrix}.$$

and

$$\Delta_r = \begin{pmatrix} t_{k+1} - \tilde{t}_{k+1} & 0 & \cdots & 0 \\ t_{k+2} - \tilde{t}_{k+2} & t_{k+1} - \tilde{t}_{k+1} & & \vdots \\ \vdots & & & \vdots \\ t_{k+r} - \tilde{t}_{k+r} & \cdots & \cdots & t_{k+1} - \tilde{t}_{k+1} \end{pmatrix}.$$

We can compute  $G_{k+r}$  by using the Gohberg-Semencul formula for  $\tilde{T}_{k+r}^{-1}$  in terms of  $\tilde{x}^{k+r} = x$  and  $\tilde{y}^{k+r} = y$  in the form

$$(10) \quad (I_r \quad 0) \tilde{T}_{k+r}^{-1} \begin{pmatrix} 0 \\ I_r \end{pmatrix} = L_1(1:r, 1:r) U_1(1:r, k+1:k+r) - L_2(1:r, 1:r) U_2(1:r, k+1:k+r)$$

in  $O(r^3)$  operations. Then we check the singular values of  $G_{k+r}$ . If  $G_{k+r}$  is near-singular we choose a new  $\tilde{t}_{k+r+1}$  and compute  $(\tilde{x}^{k+r+1}, \tilde{y}^{k+r+1})$ . We always try to choose the new element equals the original  $t_{k+r+1}$ , and if this is not possible we force  $\gamma^{k+r+1}$  or  $1 - \omega_0^{k+r+1} z_0^{k+r+1}$  to be 1.

If the singular values of  $G_{k+r}$  are not too small we have to jump back to the regular original matrix  $T_{k+r}$ . This matrix has near singular submatrices, and therefore we use  $x$  and  $z$  instead of  $x$  and  $y$  in order to characterize  $T_{k+r}^{-1}$ . First, we compute  $\tilde{z}^{k+r}$  using (6) which needs in addition one inner product. Then we have to compute with  $m = k+r$

$$T_m^{-1}(e_0, \tau_-^m) = (\tilde{x}^m, \tilde{z}^m) - \tilde{T}_m^{-1} \begin{pmatrix} 0 \\ I_r \end{pmatrix} G_m^{-1} (\Delta_r \quad 0) (\tilde{x}^m, \tilde{z}^m) = (\tilde{x}^m, \tilde{z}^m) - \tilde{T}_m^{-1} \left( \begin{pmatrix} 0 \\ u \end{pmatrix}, \begin{pmatrix} 0 \\ v \end{pmatrix} \right).$$

Now, any solution of  $\tilde{T}_{k+r} a = b$  can be computed recursively by (see [9,p.60])

$$a^j = \begin{pmatrix} a^{j-1} \\ 0 \end{pmatrix} - J \tilde{y}^j (\tilde{\tau}_+^m a^{j-1} - b_j).$$

In our case, the vector  $b$  has  $k+1$  leading zeros and therefore the first iterates are zero. This leads to the recursions

$$(11) \quad \phi^k = 0_{k+1}, \quad \phi^{k+j} = \begin{pmatrix} \phi^{k+j-1} \\ 0 \end{pmatrix} - J \tilde{y}^{k+j} (\tilde{\tau}_+^{k+j} \phi^{k+j-1} - u(j))$$

$$(12) \quad \psi^k = 0_{k+1}, \quad \psi^{k+j} = \begin{pmatrix} \psi^{k+j-1} \\ 0 \end{pmatrix} - J \tilde{y}^{k+j} (\tilde{\tau}_+^{k+j} \psi^{k+j-1} - v(j))$$

for  $j = 1, \dots, r$ . Then we get the new generating vectors for  $T_{k+r}$  by

$$(13) \quad (x^{k+r}, z^{k+r}) = (\tilde{x}^{k+r}, \tilde{z}^{k+r}) - (\phi^{k+r}, \psi^{k+r}).$$

Now, we have to compute a new pair  $(x, y)$  on level  $(k + r + 1)$  from  $(x, z)$ . For  $x^{k+r+1}$  we use (3), which also shows whether  $T_{k+r+1}$  is singular or not. The vector  $y^{k+r+1}$  results from (6) without using additional inner products. Moreover, from (7) we get  $w_0^{k+r}$  without using further inner products. Hence, the algorithm takes in every step two inner products plus two SAXPY's of size  $k$  for computing the generating vectors, and in the case of breakdown we need in addition 2 inner products plus 2 SAXPY's of size  $k + 1, k + 2, \dots, k + r$  for updating  $\phi$  and  $\psi$ . If  $T_{k+r+1}$  again is near singular then we again have a breakdown, and we compute  $\tilde{x}^{k+r+1}$  and  $\tilde{y}^{k+r+1}$  using (3), (6), and (7), after perturbing  $T_{k+r+1}$ , in order to get a regular matrix  $\tilde{T}_{k+r+1}$ .

If  $t_0 = 0$ , then the algorithm starts with a breakdown. Therefore, in this case we set  $x = z = []$ ,  $w_z - t_0 = 0$ , and begin with the breakdown case.

### BASIC ALGORITHM:

INPUT:  $t_{-n}, \dots, t_0, \dots, t_n, ep1, ep2$

OUTPUT:  $(x, y)$  or  $(x, z)$ , generating vectors for Gohberg-Semencul formulas

Initialisation:

If  $|t_0| > ep1$ : Set  $x = y = 1/t_0$ ,  $w_0 = t_1x$ ,  $z_0 = t_{-1}y$ ,  $m = 1 - w_0z_0$ ,  $k = 0$

If  $|t_0| \leq ep1$ : Set  $x = y = []$ ,  $m = 0$ ,  $k = -1$

Repeat until ready:

a) If  $|m| > ep1$ : (Regular step)

compute  $(x, y)$  according to (1-2) or  $(x, y)$  according to (3)(6);

compute  $w_0, z_0$ ,  $m = 1 - w_0z_0$ ; set  $k = k + 1$ ; if  $k = n - 1$ : ready;

b) If  $|m| \leq ep1$ : (Irregular step)

$\alpha 1$ ) if  $k = -1$ : Set  $\tilde{t}_0 = \tilde{y} = \tilde{x} = -1$ ,

$\alpha 2$ ) If  $k \geq 0$ : Set  $\tilde{t} = (t_0, \dots, t_k)$ ;

compute  $(\tilde{x}, \tilde{y}_1)$  (1-2) with  $\tilde{t}_{k+1} = \frac{m-1}{x_0z_0} + t_{k+1}$ ,  $w_0 = w_0 + x_0(\tilde{t}_{k+1} - t_{k+1})$ ;

or

compute  $(\tilde{x}, \tilde{y}_1)$  (3),(6) with  $\tilde{t}_{k+1} = \frac{1-m}{z_0} + t_{k+1}$ ,  $w_0 = w_0 + x_0(\tilde{t}_{k+1} - t_{k+1})$ ;

$\beta$ ) Repeat until regular:

$r = 2$ , compute  $w_0, z_0, m = 1 - w_0z_0$ ;

1) If  $|m| > ep1$ :  $\tilde{t}_{k+r} = t_{k+r}$ , compute  $(\tilde{x}, \tilde{y}_r)$  with (1-2);

2) If  $|m| \leq ep1$ : compute  $(\tilde{x}, \tilde{y}_r)$  using (1-2) with  $\tilde{t}_{k+r} = \frac{m-1}{x_0z_0} + t_{k+r}$ ,

$m = 1, w_0 = w_0 + \tilde{x}_0(\tilde{t}_{k+r} - t_{k+r})$ ;

$\gamma$ ) Compute  $G$  from  $(\tilde{x}, \tilde{y}_r)$  with (10);

$\delta$ ) 1) If  $\min(\text{svd}(G)) \leq ep2$ :  $r = r + 1$ , singular;

2) If  $\min(\text{svd}(G)) > ep2$ : compute  $\tilde{z}$  with (6);

$(u, v) = G^{-1} \Delta_r((\tilde{x}, \tilde{z})(1 : r))$ ;

$\phi_i = \psi_i = 0_{k+1}$ ; compute  $\phi_i, \psi_i$  with (11-12) using  $\tilde{y}_1, \dots, \tilde{y}_r$ ,

and  $(x, z)$  with (13)

Set  $k = k + r$ , if  $k = n - 1$ : ready;

$m = (t_{k+2}, \dots, t_0) \begin{pmatrix} z \\ -1 \end{pmatrix}$ ,  $w_0 = \frac{1+mz_0}{z_0}$ ; regular.

Note, that  $m$  is defined in two different ways,  $m = 1 - w_0z_0$  or  $m = w_z - t_0$ , dependent on what formula is used for computing  $(x, y)$ .

The numerical examples are taken from [4]. We consider Toeplitz matrices of the form  $T = T^0 + \epsilon T^1$  where the entries of  $T^1$  are randomly distributed in  $[0, 1]$  and  $T^0$  is given as



the coefficients of the Taylor expansion of

$$f(z) = 18 + \frac{1}{1-z} - \frac{3}{1-z^3} + \frac{6}{1-z^6} - \frac{24}{1-z^{24}} + \frac{48}{1-z^{48}} - \frac{96}{1-z^{96}}.$$

The submatrices of dimension 51, ..., 57 are exactly singular, while all other submatrices are well-conditioned. The second example is given by  $T_n(\epsilon)$  with  $t_0 = \epsilon$ ,  $t_j = 1/2^j$ ,  $j = 1, \dots, n$ . Here for  $\epsilon = 0$  every third submatrix is exactly singular.

We apply our look-ahead algorithm for different values of  $\epsilon$  and display the relative error  $r = \|x - T^{-1}e_0\|/\|x\|$ .  $ep1$  and  $ep2$  are chosen to be 0.01.

$\epsilon$	1.0e-02	1.0e-04	1.0e-06	1.0e-08	1.0e-10	1.0e-12	1.0e-14	0.0e+00
r	1.2e-13	6.6e-13	4.4e-13	7.5e-13	3.9e-13	5.5e-13	3.6e-13	7.2e-13

Table 1. Relative error for example 1 with different values of  $\epsilon$

For  $\epsilon = 0.01$  there occurred no breakdown in view of  $ep1 = 0.01$ , but for all other values of  $\epsilon$  the algorithm made a jump from 51 to 58.

n+1	15	30	60	120	240	480
r	3.11e-15	3.94e-15	1.08e-14	9.87e-14	4.19e-13	1.18e-12

Table 2. Relative error for example 2 with  $\epsilon = 10e - 14$

Here the algorithm had a jump of height 2 every third step. If we compare this results with the relative errors for the look-ahead Bareiss algorithm given in [4], we see that our relative errors are only slightly larger, but our computational effort is much smaller.

The algorithm here is given in its fastest version. We can consider different variations of the basic algorithm, e.g. for the symmetric case, which can be found in [10]. For other look-ahead Toeplitz solvers see [3,4,5,8].

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